JAVACC PARSER OPTIONS

The following JavaCC options may be useful for debugging your work: DEBUG_LOOKAHEAD =true; DEBUG_PARSER=true;

RUDIMENTS

- Upper and lower case reversed for terminals and non-terminals.
- All non-terminals are function calls.
- After Token definitions:

```
void non-terminal() :
    { declarations }
    {     prod
          prod
          prod
          prod
     }
```

- Tokens: either <NAME> or "actual string" allowed
- Shorthands: |*+? allowed (x)? = [x]
- ε productions:

```
{} /* nothing */
```

- Or-ed productions are tried in the order presented
- Example:

```
IF_STAT \rightarrow "if" COND "then" STAT "else" STAT "end"
IF_STAT \rightarrow "if" COND "then" STAT "end"
```

```
void if_stat() :
{}
{ "if" cond() "then" stat() "else" stat() "end"
| "if" cond() "then" stat() "end"
}
```

LL ISSUES

Global Lookaheads

- Default: JavaCC assumes language is LL(1)
- Can be made LL(k) by setting global LOOKAHEAD(k) at top of file - Unacceptable as previously discussed

Local Lookaheads

• Can use local lookahead specific to a specific point in a specific production, called a decision point.

```
void S() :
{}
{ "a" "b" "c"
| "a" "d" "c"
}
```

Decision point right before first "a" → replace by:

```
void S() :
{}
{ LOOKAHEAD(2) "a" "b" "c"
    "a" "d" "c"
}
```

• Second Example: void S() :

{}
{ "a" "b" "0"
| "a" "b" "1"
}

```
Solution 1 – no factoring
```

```
void S() :
{}
{ LOOKAHEAD(3)"a" "b" "0"
| "a" "b" "1"
}
```

Solution 2 – partial factoring

```
void S() :
{}
{ "a" (LOOKAHEAD(2) "b" "0" | "b" "1")
}
```

```
Solution 3 - full factoring
void S() :
{}
{ "a" "b"("0"|"1")
}
```

• Compare and discuss backtracking.

Syntactic Lookaheads

• Example:

```
void S() :
{}
{ ("a")+ "0"
| ("a" | "b")+ "1"
}
```

Don't know how many letters to look ahead

• Solution:

```
void S() :
{}
{ LOOKAHEAD(("a")+ "0") ("a")+ "0"
| ("a" | "b")+ "1"
}
```

- How much can it lookahead?
 - Possibly the entire program
 - − VERY COSTLY \rightarrow AVOID!!!
 - Very few non-terminals in the assignment need them.
- In reality your program would probably look like this:

```
void S() :
{ }
    lots of _as_then_0()
{
   as and \overline{bs}() "1"
}
void lots of as then 0 () :
{ }
     ("a") + "0"}
{
void as and bs() :
{ }
     ("a" | "b")<sup>+</sup>
{
}
```

You may not notice until JavaCC tells you about a choice conflict in S. \rightarrow resolution:

```
void S() :
{}
{ LOOKAHEAD(lots_of_as_then_0 ()) lots_of_as_then_0 ()
| as_and_bs() "1"
}
```

- Where to put the syntactic lookahead?
 - where you expect the shortest matching string, or the most likely string to be matched correctly so there is no need to backtrack.

Lookahead-only Productions

• Example

```
void declaration() :
{ }
  LOOKAHEAD(fn declaration()) fn declaration()
{
    fn definition()
other declaration()
}
void fn definition():
{ }
   type() <IDENTIFIER> "(" parameters() ")" "{" body() "}"
{
}
void fn declaration():
{ }
    type() <IDENTIFIER> "(" parameters() ")" ":" package()
{
";"
}
```

Don't want to read entire definition or declaration to decide which it is. \rightarrow define a production simply for looking-ahead:

```
void fn_decl_lookahead():
{}
{ type() <IDENTIFIER> "(" parameters() ")" ":"
}
void declaration() :
{}
{ LOOKAHEAD(fn_decl_lookahead()) fn_declaration()
| fn_definition()
| other_declaration()
}
```